# Elementary Numbah Theory

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## 1 Axioms

Axiom 1.  $1 \in \mathbb{N}$ 

Axiom 2.  $\forall x \in \mathbb{N}, \exists x', \text{ called the successor of } x$ 

Axiom 3. 1 is not the successor of any natural number

Axiom 4.  $x' = y' \Rightarrow x = y$ 

Axiom 5. : Induction Axiom Let  $S \subseteq \mathbb{N}$  such that • (1)  $1 \in S$ • (2)  $x \in S \Rightarrow x' \in S$ 

Then  $S = \mathbb{N}$ 

Axiom 6. : Well-Ordering Axiom  $\forall (S \subseteq \mathbb{N} \land S \neq \emptyset), S \text{ contains a least element.}$ That is,  $\forall (b \in S), (\exists a \in S) : a \leq b$ 

# 2 Postulates on $\mathbb{Z}$

**Postulate 1.** Reflexivity of Equality  $a \in \mathbb{Z} \Rightarrow a = a$ 

**Postulate 2.** Symmetry of Equality  $a, b \in \mathbb{Z} \land a = b \Rightarrow b = a$ 

**Postulate 3.** Transitivity of Equality  $a, b, c \in \mathbb{Z} \land a = b \land b = c \Rightarrow a = c$ 

**Postulate 4.** Transitivity of Inequality  $a, b, c \in \mathbb{Z} \land a < b \land b < c \Rightarrow a < c$ 

**Postulate 5.** Trichotomy  $a, b \in \mathbb{Z} \Rightarrow$  Exactly one of the following is true: (1) a < b, (2) a = b or (3) a > b

Postulate 6. Binary operations

Multiplication Addition .  $(ab) \in \mathbb{Z}$ Closure  $(a+b) \in \mathbb{Z}$ Equality  $a = b \Rightarrow a + c = b + c$  $a = b \Rightarrow a \cdot c = b \cdot c$ Associativity (a+b) + c = a + (b+c) (ab)c = a(bc)Identity a + 0 = 0 + a = a $a \cdot 1 = 1 \cdot a = a$ Commutativity a+b=b+a $a \cdot b = b \cdot a$ Inverse a + (-a) = 0 $\{1, -1\}$ Transitivity of Inequality  $a < b \Leftrightarrow a + c < b + c$  $a < b \Leftrightarrow a \cdot |c| < b \cdot |c|$  $a \cdot (b+c) = ab + ac$ Distributivity

 $\mathbb{Z}^+$  is an abelian group and an infinite cyclic group

 $\mathbb{Z}^*$  is a commutative monoid

# 3 Divisibility

**Definition 3.1.** Let  $a, b \in \mathbb{Z}$  $\exists (k \in \mathbb{Z}) : b = ak \Rightarrow a \mid b$  $\coloneqq a$  divides b $\coloneqq a$  is a divisor of b

**Properties.** 

**Property 3.1.** 0 is not a divisor of any integer except 0, since  $\neg \exists k \neq 0$ :  $0 \cdot k \neq 0$ 

**Property 3.2.**  $a \mid 0$  since  $0 = 0 \cdot a$ 

**Property 3.3.**  $1 \mid a \text{ since } a = 1 \cdot a$ 

**Property 3.4.**  $a \mid a$  since  $a = a \cdot 1$ 

**Property 3.5.**  $a \mid b \land b \neq 0 \Rightarrow |a| \leq |b|$ 

P	Proof		
1.	$b = ak$ for some $k \in \mathbb{Z}$	(Definition 3.1)	
2.	$k \neq 0$ since $b \neq 0$	(Premise)	
3.	$ k  \ge 1$	(2)	
4.	$ b = ak = a \cdot  k \geq  a \cdot 1$	(3, Trans. Ineq.)	

**Property 3.6.** Closure under multiplication  $d \mid a \Rightarrow d \mid ab$ 

Proof1dk = a2 $dk \cdot b = a \cdot b$ 2 $d \cdot (kb) = a \cdot b$ 4 $d \mid ab$ 4 $d \mid ab$ Converse is not necessarily true.

**Property 3.7.** Transitivity  $a \mid b \land b \mid c \Rightarrow a \mid c$ 

	Proof		
1	ak = b for some $k$	(Definition 3.1)	
2	$b = ak \mid c$	(Premise)	
3	akk' = c for some $k'$	(Definition $3.1$ )	
4	a(kk') = c	(Assoc. Mult.)	
5	$a \mid c$	(Definition $3.1$ )	

Property 3.8. Equality

 $a \mid b \Leftrightarrow a \cdot c \mid b \cdot c$ 

 $\begin{array}{l} a \mid b \Rightarrow ak = b \Rightarrow (ac)k = bc \Rightarrow ac \mid bc \\ ac \mid bc \Rightarrow ack = bc \Rightarrow ak = b \Rightarrow a \mid b \end{array}$ 

**Property 3.9.**  $a \mid b \land b \mid a \Rightarrow |a| = |b|$ 

Proof

 $\begin{array}{l} a \mid b, \text{ so } ak = b \text{ for some } k \in \mathbb{Z} \\ b \mid a, \text{ so } bk' = a \text{ for some } k' \in \mathbb{Z} \\ akk' = a \Leftrightarrow (k, -k) \in \{(1, -1), (-1, 1)\} \end{array}$ 

# 4 Common Divisor

**Definition 4.1.** Let  $d \mid a$  and  $d \mid b, d \in \mathbb{Z}$ := d is a **common divisor** of a and b

**Definition 4.2.** A linear combination of  $a, b \in \mathbb{Z}$  is any integer of the form  $ra + sb, r, s \in \mathbb{Z}$ 

**Theorem 4.1.** Linear Combination Let  $(a, b, d, r, s) \in \mathbb{Z}$ . Then  $d \mid a \land d \mid b \Rightarrow d \mid (ra + sb)$ 

Proof $d \mid a$ Premise $d \mid b$ Premise $(1) \quad \exists (e \in \mathbb{Z}) : a = d \cdot e$ (Definition 3.1) $(2) \quad \exists (f \in \mathbb{Z}) : b = d \cdot f$ (Definition 3.1) $(3) \quad ra + sb = rde + sdf = d(re + sf)$ (1, 2) $(4) \quad d \mid (ra + sb)$ (3), Definition 3.1Note that the converse is not necessarily true.

**Corollary 4.1.1.**  $d \mid a \land d \mid b \Rightarrow d \mid (a+b)$ Set r = 1, s = 1

**Corollary 4.1.2.**  $d \mid a \land d \mid b \Rightarrow d \mid (a - b)$ Set r = 1, s = -1 Corollary 4.1.3.  $d \mid a \Rightarrow d \mid ra$ Set r = 1, s = 0. Also, see (3.6)

**Lemma 4.2.** For a, b not both 0, there is a least positive linear combination of a and b.

*Proof.* WLOG, assume  $a \neq 0$ . Let  $S = \{x : x = (r_0 a + s_0 b) \ \forall (r_0, s_0 \in \mathbb{Z}) \}$ . Then  $a \in S$  for a > 0,  $r_0 = 1$  and  $-a \in S$  for a < 0,  $r_0 = 1$ . Therefore,  $S \neq \emptyset$ . 

**Lemma 4.3.** For a, b not both 0, the least positive linear combination of a and b is a common divisor of a and b.

Proof

FTOOJ		
(1)	Let $d$ be the least positive linear combination of $a$ and $b$	(Lemma 4.2)
(2)	Write $a = qd + r, \ 0 \le r < d$	(Division Algorithm)
(3)	$r = a - qd = a - q(r_0a + s_0b) = (1 - qr_0)a + (-qs_0)b$	(2)
(4)	$r$ is also a linear combination of $a$ and $b$ and $r \geq 0$	(2,3)
(5)	If $r > 0$ then (1) is contradicted. Therefore, $r = 0$	(1, 4)
(6)	$a = qd + 0$ , Hence $a = qd$ and $d \mid a$	(2)
(7)	Repeat $(2)$ - $(6)$ with b to complete the proof	

#### $\mathbf{5}$ Greatest Common Divisor

**Definition 5.1.** For a, b not both 0, there is a greatest common divisor of a and b

*Proof.* WLOG, assume  $a \neq 0$ . Let  $S = \{x : x \mid a \land x \mid b\}$ . Then: Existence  $1 \in S \Rightarrow S \neq \emptyset$ (Property 3.3) Upper bound  $x \in S \Rightarrow x \mid a \Rightarrow x \leq |a|$  (Property 3.5)

**Properties.** 

**Property 5.1.** For a, b not both 0,  $gcd(a, b) \ge 1$  since  $1 \in S$ 

**Property 5.2.** gcd(a, 0) = a since  $a \mid 0$  and  $a \mid a$ 

**Property 5.3.**  $c \mid a \land c \mid b \Rightarrow c \mid qcd(a, b)$ . This follows directly from Theorem 4.1.

**Property 5.4.**  $gcd(ac, bc) = c \cdot gcd(a, b)$ 

Proof Let d = gcd(a, b)Let d' = gcd(ac, bc)(1) Then  $d \mid a \land d \mid b$  from Definition 4.1 (2) And  $dc \mid ac \wedge dc \mid bc$  from Property 3.8 (3) So  $dc \mid d'$  from Property 5.3 (4) d = ra + sb for some r and s from (1) (5) Then dc = rac + sbc (Multiplicative Equality) (6)  $d' \mid ac \wedge d' \mid bc$  by Definition 5.1

(7)  $d' \mid (rac + sbc) \Rightarrow d' \mid dc$  from Theorem 4.1, (5)

(8) From (3),  $dc \leq d'$ , and from (7),  $d' \geq dc$ 

(9) Therefore  $d' = dc \square$ 

#### Theorem 5.1. Bezout's Identity

gcd(a, b) is the least positive linear combination of a and b

Proof

Let *d* be the least positive linear combination of *a* and *b*. Then *d* | *a* and *d* | *b* from Lemma 4.3 Let *c* | *a* and *c* | *b* for some  $c \in \mathbb{Z}$ Since *d* is a linear combination of *a* and *b*, *c* | *d* And  $c \leq d$  from Property 3.5 All common divisors *c*, of (*a*, *b*) are  $\leq d \Rightarrow gcd(a, b) = d$ 

**Corollary 5.1.1.**  $c \mid gcd(a, b) \Leftrightarrow c \mid a \land c \mid b$ 

 $Proof \leftarrow$  is restatement of Property 5.3

 $\begin{array}{l} Proof \Rightarrow \\ c \mid gcd(a,b) \text{ by Premise} \\ gcd(a,b) \mid a \text{ by Definition 5.1} \\ c \mid a \text{ by Property 3.7} \\ c \mid b \text{ can be shown by analogous derivation. } \end{array}$ 

**Corollary 5.1.2.**  $\forall k \in \mathbb{Z}, \exists (r, s) \in \mathbb{Z} : k \cdot gcd(a, b) = ra + sb$ All multiples of gcd(a, b) are a linear combination of a and b

 $\begin{array}{ll} Proof\\ gcd(a,b) = r_0 a + s_0 b & \text{Theorem 5.1}\\ k \cdot gcd(a,b) = (r_0 k)a + (s_0 k)b & \Box \end{array}$ 

**Corollary 5.1.3.**  $\forall (a, b, r, s \in \mathbb{Z}) : gcd(a, b) \mid (ra + sb)$ All linear combinations of a and b are a multiple of gcd(a, b)

 $\begin{array}{ll} Proof\\ \text{Let }g=gcd(a,b)\\ \text{Let }(r,\,s)\in\mathbb{Z} \text{ be arbitrary integers}\\ k\cdot g=a & \text{Definition 5.1}\\ k'\cdot g=b & \text{Definition 5.1}\\ ra+sb=rkg+sk'g=(rk+sk')g & \text{Substituting}\\ g\mid (ra+sb) & \Box \end{array}$ 

Corollary 5.1.4.  $gcd(ac, bc) = c \cdot gcd(a, b)$ 

Proof

	Let $d = gcd(a, b)$	
1	$d \mid a \wedge d \mid b$	Definition 4.1
2	$dc \mid ac \wedge dc \mid bc$	Property 3.8
	Let $d' = gcd(ac, bc)$	
3	Then $dc \mid d'$	(2), Property 5.3
4	$d = r_0 a + s_0 b$	Theorem 5.1
5	$dc = r_0 ac + s_0 bc$ , which is a linear combination of $ac, bc$	
6	$d' \mid dc$	(5), Corollary 5.1.3
7	d' = dc	$(3), (6) \square$

Corollary 5.1.5.  $gcd(a, bc) \mid (gcd(a, b) \cdot gcd(a, c))$ 

 $\begin{array}{ll} Proof \\ \text{Let } gcd(a,b) = r_0a + s_0b & \text{Theorem 5.1} \\ \text{Let } gcd(a,c) = r_1a + s_1c & \text{Theorem 5.1} \\ \text{Then } (gcd(a,b) \cdot gcd(a,c) = r_0ar_1a + r_0as_1c + s_0br_1a + s_0bs_1c \\ = (r_0r_1a + r_0s_1c + s_0s_1b)a + (s_0s_1)bc & \text{which is a linear combination of } a, bc \\ \text{And which } gcd(a,bc) \text{ is thus a divisor of by Corollary 5.1.3} & \Box \end{array}$ 

**Corollary 5.1.6.** gcd(a + bc) = gcd(a, b) for any  $c \in \mathbb{Z}$ 

Proof			
	Let $d = gcd(a, b)$ and $d' = gcd(a + bc, b)$		
1	d' = r(a+bc) + sb = ra + (rc+s)b	Definition 4.1	
2	$d \mid a \wedge d \mid b$	Definition 4.1	
3	$d \mid d'$	(2), Theorem 4.1	
4	$d' \mid b \Rightarrow d' \mid bc$	Property 3.6	
5	$d' \mid (a + bc) \land d' \mid bc \Rightarrow d' \mid (a + bc - bc) \Rightarrow d' \mid a$	Corollary 4.1.2	
6	$d' \mid a \wedge d' \mid b \Rightarrow d' \mid d$	Property 5.3	
7	$d \mid d' \wedge d' \mid d \Rightarrow d = d'$	$(3), (6) \square$	

#### Lemma 5.2.

For a, b not both 0, write a = bq + r. Then gcd(a, b) = gcd(b, r)

 $\begin{array}{l} Proof\\ \text{Let } d = gcd(a,b)\\ \text{Let } d' = gcd(b,r)\\ 1 \quad d = sa + tb = s(bq+r) + tb = (sq+t)b + sr \text{ for some } s,t \in \mathbb{Z}\\ 2 \quad \text{Then } d' \mid d\\ 3 \quad d' = s'b + t'r = s'b + t'(a - bq) = t'a + (s - qt')b \text{ for some } s',t' \in \mathbb{Z}\\ 4 \quad \text{Then } d \mid d'\\ \text{From } (2) \text{ and } (4), d = d' \end{array}$ 

# 6 Euclidean GCD Algorithm

Consider the following sequence of divisions:

 $a = bq_0 + r_0$  $0 \le r_0 < b$  $gcd(a,b) = gcd(b,r_0)$  $\begin{array}{ll} 0 \leq r_1 < r_0 & gcd(b.r_0) = gcd(r_0, r_1) \\ 0 \leq r_2 < r_1 & gcd(r_0, r_1) = gcd(r_1, r_2) \\ 0 \leq r_3 < r_2 & gcd(r_1, r_2) = gcd(r_2, r_3) \end{array}$  $b = q_1 r_0 + r_1$  $r_0 = q_2 r_1 + r_2$  $r_1 = q_3 r_2 + r_3$ ... ... ••• ... ••• ...  $r_{n-2} = q_n r_{n-1} + r_n \quad 0 \leq r_n \leq r_{n-1} \quad \gcd(r_{n-2}, r_{n-1}) = \gcd(r_{n-1}, r_n)$  $gcd(r_{n-1}, r_n) = gcd(r_n, 0) = r_n$  $r_{n-1} = q_{n+1}r_n + 0$ Note that the sequence  $r_0, r_1, r_2...r_n$  is strictly decreasing.

Therefore, it will eventually yield 0.

Let  $r_n$  be the last non-zero remainder. Now:

 $r_n \mid r_{n-1}$  from the last term

So  $r_n \mid r_{n-2}$  from the term above

Proceeding similarly,  $r_n \mid b$  and  $r_n \mid a$ , so  $r_n$  is a common divisor of a and bLet d be an arbitrary common divisor of a, b. Then:

 $\begin{array}{l} d \mid (a - bq_0) \Rightarrow d \mid r_0 \\ d \mid (b - q_1 r_0) \Rightarrow d \mid r_1 \\ d \mid (r_0 - q_2 r_1) \Rightarrow d \mid r_2 \\ \dots \end{array}$ 

 $d \mid (r_{n-2} - q_n r_{n-1}) \Rightarrow d \mid r_n$ 

Since an arbitrary common divisor of (a, b) divides  $r_n$ ,  $r_n = gcd(a, b)$  (Property 5.3)

 $r_n = gcd(a, b)$  can also be observed by noting the sequence in the right hand column, which follows from Lemma 5.2.

# 7 Coprimality

**Definition 7.1.** Let gcd(a, b) = 1:= a and b are **coprime** := a and b are **relatively prime** := 1 is the only common divisor of a and b

**Property 7.1.**  $gcd(a,b) = 1 \Leftrightarrow ra + sb = 1$  for some  $r, s \in \mathbb{Z}$ This follows directly from Theorem 5.1

**Proposition 7.1.**  $\frac{a}{qcd(a,b)}$  and  $\frac{b}{qcd(a,b)}$  are coprime

Proof

Let d = gcd(a, b)

- 1  $d \mid a \Rightarrow dk = a$  for some  $k \in \mathbb{Z}$
- 2  $d \mid b \Rightarrow dk' = b$  for some  $k' \in \mathbb{Z}$
- 3  $\frac{a}{d} = k, \ \frac{b}{d} = k'$

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- 4 Suppose k and k' have a common divisor, d'
  - Then  $k = d'm, \ k' = d'm'$  for some  $m, m' \in \mathbb{Z}$
- 6 So a = dd'm and b = dd'm', which means dd' is a common divisor of a, b

(1), (2)

(4)

- 7 But d is the greatest common divisor of a, b, so d' = 1
- 8  $d' = 1, \Rightarrow k, k'$  are relatively prime

Theorem 7.1. Generalized Euclid's Lemma

 $a \mid bc \wedge gcd(a,b) = 1 \Rightarrow a \mid c$ 

ProofDefinition 3.11ak = bc for some  $k \in \mathbb{Z}$ Definition 3.121 = ra + sb for some  $r, s \in \mathbb{Z}$ Property 7.13c = rac + sbc = rac + sak = a(rc + sk)(1), (2)4 $a \mid c$ (3)  $\Box$ 

**Corollary 7.1.1.** Euclid's Lemma For any prime,  $p, p \mid bc \Rightarrow p \mid b \lor p \mid c$ 

Proof WLOG, assume  $p \nmid b$ Then gcd(p,b) = 1 = rp + sb for some  $r, s \in \mathbb{Z}$ Multiplying by c, c = (rc)p + (s)bcSince  $p \mid (rc)p \land p \mid (s)bc, p \mid c$ The derivation for  $p \mid b$  is analogous

# 8 Linear Diophantine Equation

#### Definition 8.1.

A Linear Diophantine Equation in 2 variables is an equation of the form ax + by = c

**Property 8.1.** ax + by = c is solvable  $\Leftrightarrow gcd(a, b) \mid c$ 

ProofLet ax + by = c have a solutionPremiseThen  $gcd(a, b) \mid c$ Corollary 5.1.3

**Theorem 8.1.** If  $(x_0, y_0)$  is a solution of ax + by = c, then all solutions are given by  $\left(x_0 + \frac{b}{gcd(a,b)}k, y_0 - \frac{a}{gcd(a,b)}k\right)$ 

 $Proof \Rightarrow$ 

Let d = gcd(a, b)c = dkProperty 8.1 d = ra + sb for some  $r, s \in \mathbb{Z}$ Theorem 5.1 c = a(rk) + b(sk)So this equation has the solution  $(x_0 = rk, y_0 = sk)$ Substituting for arbitrary x, y:  $ax + by = a\left(x_0 + \frac{b}{d}k\right) + b\left(y_0 - \frac{a}{d}k\right)$  $= (ax_0 + by_0) + \frac{abk}{d} - \frac{abk}{d} = (ax_0 + by_0)$  $Proof \Leftarrow$ Let  $x_0, y_0$  and  $(x_1, y_1)$  be solutions of ax + by = cThen  $c = ax_0 + by_0 = ax_1 + by_1$  $a(x_1 - x_0) = b(y_0 - y_1)$  $a(x_1 - x_0) = b(y_0 - y_1)$   $\frac{a}{d}(x_1 - x_0) = \frac{b}{d}(y_0 - y_1)$   $gcd(\frac{a}{d}, \frac{b}{d}) = 1$ Since  $\frac{b}{d} \nmid \frac{a}{d}, \frac{b}{d} \mid (x_1 - x_0)$ Hence,  $x_1 - x_0 = \frac{b}{d}k$  for some  $k \in \mathbb{Z}$ And  $x_1 = x_0 + \frac{b}{d}k$   $a\frac{b}{d}k = b(y_0 - y_1)$   $\frac{a}{d}k = (y_0 - y_1)$   $y_1 = y_0 - \frac{a}{d}k$ Proposition 7.1 

**Corollary 8.1.1.** For gcd(a, b) = 1, all solutions of ax + by = 1 are given by  $(x_0 + bk, y_0 - ak) \quad \forall k \in \mathbb{Z}$ , where  $(x_0, y_0)$  is one solution.

Substitute 1 for d in Proof above.

## 9 Congruence

**Definition 9.1.** If  $m \mid (a - b)$ :

 $\coloneqq a \equiv b \! \mod m$ 

 $\coloneqq a \text{ is congruent to } b \mod m$ 

:= a and b are in the same congruence class